

Difference Distribution Table of a Regular Substitution Box

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March 26, 1997

This short paper reports an interesting property of the difference distribution table of an S-box or substitution box, which has been discovered by the authors while studying relationships between differential and other cryptographic characteristics of an S-box. Namely, an $n \times m$ S-box is regular if and only if the sum of the entries in a column in the difference distribution table of the S-box is 2^{2n-m} .

Denote by V_n the vector space of n tuples of elements from $GF(2)$. An $n \times m$ S-box is a mapping from V_n to V_m , i.e., $F = (f_1, \dots, f_m)$, where n and m are integers with $n \geq m \geq 1$ and each component function f_j is a function from V_n to $GF(2)$ (or on V_n for short).

The *Sylvester-Hadamard matrix* (or *Walsh-Hadamard matrix*) of order 2^n , denoted by H_n , is generated by the recursive relation

$$H_n = \begin{bmatrix} H_{n-1} & H_{n-1} \\ H_{n-1} & -H_{n-1} \end{bmatrix}, \quad n = 1, 2, \dots, \quad H_0 = 1.$$

Each row (column) of H_n is a linear sequence of length 2^n .

In cryptography we are mainly concerned with *regular* S-boxes. An S-box $F = (f_1, \dots, f_m)$ is said to be regular if $F(x)$ runs through each vector in V_m 2^{n-m} times while x runs through V_n once. It is well-known that a regular S-box can be characterized by the balance of the linear combinations of its component functions. The following is a re-statement of Corollary 7.39 of [1]:

Lemma 1 *Let $F = (f_1, \dots, f_m)$ be a mapping from V_n to V_m , where n and m are integers with $n \geq m \geq 1$ and each $f_j(x)$ is a function on V_n . Then F is regular if and only if every non-zero linear combination of f_1, \dots, f_m , $f(x) = \bigoplus_{j=1}^m c_j f_j(x)$, is balanced.*

Now we introduce three notations: $k_j(\alpha)$, $\Delta_j(\alpha)$ and η_j associated with $F = (f_1, \dots, f_m)$.

Definition 1 *Let $F = (f_1, \dots, f_m)$ be an $n \times m$ S-box, $\alpha \in V_n$, $j = 0, 1, \dots, 2^m - 1$ and $\beta_j = (b_1, \dots, b_m)$ be the vector in V_m that corresponds to the binary representation of j . In addition, set $g_j = \bigoplus_{u=1}^m b_u f_u$ be the j th linear combination of the component functions of F . Then we define*

1. $k_j(\alpha)$ as the number of times $F(x) \oplus F(x \oplus \alpha)$ runs through $\beta_j \in V_m$ while x runs through V_n once.
2. $\Delta_j(\alpha)$ as the auto-correlation of g_j with shift α .
3. η_j as the sequence of g_j .

Using the three notations we introduce three matrices in the following:

Definition 2 For $F = (f_1, \dots, f_m)$, set

$$K = \begin{bmatrix} k_0(\alpha_0) & k_1(\alpha_0) & \dots & k_{2^m-1}(\alpha_0) \\ k_0(\alpha_1) & k_1(\alpha_1) & \dots & k_{2^m-1}(\alpha_1) \\ & \vdots & & \\ k_0(\alpha_{2^n-1}) & k_1(\alpha_{2^n-1}) & \dots & k_{2^m-1}(\alpha_{2^n-1}) \end{bmatrix},$$

$$D = \begin{bmatrix} \Delta_0(\alpha_0) & \Delta_1(\alpha_0) & \dots & \Delta_{2^m-1}(\alpha_0) \\ \Delta_0(\alpha_1) & \Delta_1(\alpha_1) & \dots & \Delta_{2^m-1}(\alpha_1) \\ & \vdots & & \\ \Delta_0(\alpha_{2^n-1}) & \Delta_1(\alpha_{2^n-1}) & \dots & \Delta_{2^m-1}(\alpha_{2^n-1}) \end{bmatrix}$$

and

$$P = \begin{bmatrix} \langle \eta_0, \ell_0 \rangle^2 & \langle \eta_1, \ell_0 \rangle^2 & \dots & \langle \eta_{2^m-1}, \ell_0 \rangle^2 \\ \langle \eta_0, \ell_1 \rangle^2 & \langle \eta_1, \ell_1 \rangle^2 & \dots & \langle \eta_{2^m-1}, \ell_1 \rangle^2 \\ & \vdots & & \\ \langle \eta_0, \ell_{2^n-1} \rangle^2 & \langle \eta_1, \ell_{2^n-1} \rangle^2 & \dots & \langle \eta_{2^m-1}, \ell_{2^n-1} \rangle^2 \end{bmatrix},$$

where ℓ_i is the i th row of H_n , $i = 0, 1, \dots, 2^n - 1$. The three $2^n \times 2^m$ matrices K , D and P are called *difference distribution table*, *auto-correlation distribution table* and *correlation immunity distribution table* of the S-box F respectively.

In designing a strong S-box, many cryptographic criteria should be examined not only against component functions, but also against their linear combinations. Such criteria include those related to non-linearity, propagation characteristics and difference distribution tables. The matrix K characterizes the differential characteristics of an S-box. The matrix D indicates the auto-correlation of all linear combinations of the component functions. While the matrix P represents the inner product between the sequence of each linear combination of the component functions and each linear sequence. P is helpful in studying the correlation immunity, as well as the nonlinearity, of each linear combination of the component functions (see [2]).

As one immediately expects, the three matrices K , D and P are closely related. In particular the following result has been proven in [4]:

Theorem 1 Let $F = (f_1, \dots, f_m)$ be a mapping from V_n to V_m , where n and m are integers with $n \geq m \geq 1$ and each $f_j(x)$ is a function on V_n . Set $g_j = \bigoplus_{u=1}^m c_u f_u$ where (c_1, \dots, c_m) is the binary representation of integer j , $j = 0, 1, \dots, 2^m - 1$. Then

(i) $D = KH_m$,

(ii) $P = H_n D$,

(iii) $P = H_n K H_m$.

Using Theorem 1, we now show that a regular S-box can be completely characterized by its difference distribution table.

Corollary 1 *Let $F = (f_1, \dots, f_m)$ be a mapping from V_n to V_m , where n and m are integers with $n \geq m \geq 1$ and each f_j is a function on V_n . Then F is regular if and only if the sum of a column in the difference distribution table is 2^{2n-m} , i.e., $\sum_{\alpha \in V_n} k_i(\alpha) = 2^{2n-m}$, $i = 0, 1, \dots, 2^m - 1$.*

Proof. Compare the first rows in both sides of the formula in (iii) of Theorem 1,

$$\left(\sum_{\alpha \in V_n} k_0(\alpha), \sum_{\alpha \in V_n} k_1(\alpha), \dots, \sum_{\alpha \in V_n} k_{2^m-1}(\alpha) \right) H_m = (\langle \eta_0, \ell_0 \rangle^2, \langle \eta_1, \ell_0 \rangle^2, \dots, \langle \eta_{2^m-1}, \ell_0 \rangle^2). \quad (1)$$

Obviously, if $\sum_{\alpha \in V_n} k_i(\alpha) = 2^{2n-m}$, $i = 0, 1, \dots, 2^m - 1$. then $\langle \eta_1, \ell_0 \rangle^2 = \dots = \langle \eta_{2^m-1}, \ell_0 \rangle^2 = 0$. Note that ℓ_0 is the all-one sequence of length 2^n . Hence g_1, \dots, g_{2^m-1} are balanced, where g_1, \dots, g_{2^m-1} are defined in Theorem 1. By Lemma 1, F is regular.

Conversely, suppose F is regular. By Lemma 1, g_1, \dots, g_{2^m-1} are balanced. Hence $\langle \eta_1, \ell_0 \rangle^2 = \dots = \langle \eta_{2^m-1}, \ell_0 \rangle^2 = 0$. Note that $\langle \eta_0, \ell_0 \rangle^2 = 2^{2n}$. Rewrite (1) as

$$2^m \left(\sum_{\alpha \in V_n} k_0(\alpha), \sum_{\alpha \in V_n} k_1(\alpha), \dots, \sum_{\alpha \in V_n} k_{2^m-1}(\alpha) \right) = (2^{2n}, 0, \dots, 0) H_m.$$

This proves that $\sum_{\alpha \in V_n} k_i(\alpha) = 2^{2n-m}$, $i = 0, 1, \dots, 2^m - 1$. \square

Corollary 1 has also been obtained independently by Tapia-Recillas, Daltaubuit and Vega [3].

The following corollary shows the uniqueness of the first column of the difference distribution table of a regular mapping.

Corollary 2 *Let $F = (f_1, \dots, f_m)$ be a mapping from V_n to V_m , where n and m are integers with $n \geq m \geq 1$ and each f_j is a function on V_n . Then F is regular if and only if the sum of the leftmost column is 2^{2n-m} , i.e., $\sum_{\alpha \in V_n} k_0(\alpha) = 2^{2n-m}$.*

Proof. Multiply both sides of the equality in (iii) of Theorem 1 by e^T where, e denotes the all-one sequence of length 2^m . Hence we have

$$H_n \begin{bmatrix} k_0(\alpha_0) & k_1(\alpha_0) & \dots & k_{2^m-1}(\alpha_0) \\ k_0(\alpha_1) & k_1(\alpha_1) & \dots & k_{2^m-1}(\alpha_1) \\ \vdots & \vdots & \ddots & \vdots \\ k_0(\alpha_{2^n-1}) & k_1(\alpha_{2^n-1}) & \dots & k_{2^m-1}(\alpha_{2^n-1}) \end{bmatrix} \begin{bmatrix} 2^m \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \begin{bmatrix} \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_0 \rangle^2 \\ \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_1 \rangle^2 \\ \vdots \\ \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_{2^n-1} \rangle^2 \end{bmatrix}$$

and hence

$$2^m H_n \begin{bmatrix} k_0(\alpha_0) \\ k_0(\alpha_1) \\ \vdots \\ k_0(\alpha_{2^n-1}) \end{bmatrix} = \begin{bmatrix} \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_0 \rangle^2 \\ \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_1 \rangle^2 \\ \vdots \\ \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_{2^n-1} \rangle^2 \end{bmatrix}. \quad (2)$$

Compare the two sides of equality (2), obtaining

$$2^m \sum_{i=0}^{2^n-1} k_0(\alpha_i) = \sum_{j=0}^{2^m-1} \langle \eta_j, \ell_0 \rangle^2. \quad (3)$$

Since g_0 is the constant zero, η_0 is the all-one sequence of length 2^n and hence $\langle \eta_0, \ell_0 \rangle^2 = 2^{2n}$.

If $\sum_{\alpha \in V_n} k_0(\alpha) = 2^{2n-m}$, then from (3), $\langle \eta_1, \ell_0 \rangle^2 = \cdots = \langle \eta_{2^m-1}, \ell_0 \rangle^2 = 0$. Note that ℓ_0 is the all-one sequence of length 2^n . Hence g_1, \dots, g_{2^m-1} are balanced, where g_1, \dots, g_{2^m-1} are defined in Theorem 1. By Lemma 1, F is regular.

Conversely, if F is regular, then by Corollary 1 $\sum_{\alpha \in V_n} k_0(\alpha) = 2^{2n-m}$. \square

From Corollaries 1 and 2, we conclude that (1) an S-box is regular, (2) the sum of the first column in its difference distribution table is 2^{2n-m} , and (3) the sum of each column in the difference distribution table is 2^{2n-m} , are all equivalent statements.

References

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