# Comments on "Generating and Counting Binary Bent Sequences" 

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#### Abstract

We prove that the conjecture on bent sequences stated in "Generating and counting bent sequences", IEEE Transactions on Information Theory, IT-36 No. 5, 1990 by C.M. Adams and S.E. Tavares is false.


Let $V_{n}$ be the vector space of $n$ tuples of elements from $G F(2)$. Any map from $V_{n}$ to $G F(2), f$, can be uniquely written as a polynomial in coordinates $x_{1}, \ldots, x_{n}$ [3], [4]:

$$
f\left(x_{1}, \ldots, x_{n}\right)=\bigoplus_{v \in V_{n}} a_{v} x_{1}^{v_{1}} \cdots x_{n}^{v_{n}}
$$

where $\oplus$ denotes the boolean addition, $v=\left(v_{1}, \ldots, v_{n}\right), a_{v} \in G F(2)$. Thus we identify the function $f$ with polynomial $f$. Note that there exists a natural one to one correspondence between vectors in $V_{n}$ and integers in $\left\{0,1, \ldots, 2^{n}-1\right\}$. This allows us to order all the vectors according to their corresponding integer values. For convenience, we use $\alpha_{i}$ to denote the vector in $V_{n}$ whose integer representation is $i$. Let $f$ be a function on $V_{n}$. The ( $1,-1$ )-sequence

$$
(-1)^{f\left(\alpha_{0}\right)},(-1)^{f\left(\alpha_{1}\right)}, \ldots,(-1)^{f\left(\alpha_{2}{ }^{n}-1\right)}
$$

is called the sequence of $f(x)$.
The function $\varphi\left(x_{1}, \ldots, x_{n}\right)=a_{1} x_{1} \oplus \cdots \oplus a_{n} x_{n} \oplus \boldsymbol{c}, a_{j}, c \in G F(2)$, is called an affine function, on $V_{n}$, in particular, a linear function if $c=0$. The sequence of an affine (a linear) function is called an affine sequence (a linear sequence).

From [4] we can give an equivalent definition of bent functions. Let $\xi$ be the sequence of a function $f$ on $V_{n}$. We call $f$ a bent function and $\xi$ a bent sequence, if the scalar product
$\langle\xi, \ell\rangle= \pm 2^{\frac{1}{2} n}$ for any linear sequence $\ell$ of length $2^{n}$. Obviously bent functions on $V_{n}$ exist only for even $n$.

Consider all the $(1,-1)$-sequences of length four: $\pm(++++), \pm(++--), \pm(+-+-)$, $\pm(+--+), \pm(+++-), \pm(++-+), \pm(+-++), \pm(-+++)$, where + and - denote 1 and -1 respectively. Each of the first eight is an affine sequence. For example, $(-++-)$ is the sequence $\varphi\left(x_{1}, x_{2}\right)=1 \oplus x_{1} \oplus x_{2}:(-1)^{\varphi(0,0)}=-1,(-1)^{\varphi(0,1)}=1,(-1)^{\varphi(1,0)}=1$, $(-1)^{\varphi(1,1)}=-1$. Note that a function on $V_{2}$ is bent if and only if it is quadratic. Thus each of the second eight is a bent sequence. For example, $(--+-)$ is the sequence $g\left(x_{1}, x_{2}\right)=$ $1 \oplus x_{1} \oplus x_{1} x_{2}:(-1)^{g(0,0)}=-1,(-1)^{g(0,1)}=-1,(-1)^{g(1,0)}=1,(-1)^{g(1,1)}=-1$.

If a bent sequence of length $2^{n}$ is a concatenation of $2^{n-2}$ bent (affine) sequences of length 4 we call it bent-based (linear-based) bent sequence. Adams and Tavares conjectured that any bent sequence is either bent-based or linear-based [1]. We now prove that this conjecture is true if only if any bent function is quadratic. However there exist infinitely many bent functions with algebraic degree higher than two [2], [4]. This implies that the conjecture is not true.

The next lemma follows directly from the definition of bent-based (linear-based) bent sequences.

Lemma 1 Let $\xi$ be the sequence of a bent function, $f$, on $V_{n}(n>2)$. Then $\xi$ is bent-based (linear-based) if and only if $f\left(x_{1}^{0}, \ldots, x_{n-2}^{0}, x_{n-1}, x_{n}\right)$ is a bent (an affine) function on $V_{2}$ for any fixed vector $\left(x_{1}^{0}, \ldots, x_{n-2}^{0}\right) \in V_{n-2}$.

Note that any function on $V_{n}$ can be written as

$$
\begin{align*}
f\left(x_{1}, \ldots, x_{n}\right) & =r\left(x_{1}, \ldots, x_{n-2}\right) \oplus p\left(x_{1}, \ldots, x_{n-2}\right) x_{n-1} \oplus q\left(x_{1}, \ldots, x_{n-2}\right) x_{n} \\
& \oplus a\left(x_{1}, \ldots, x_{n-2}\right) x_{n-1} x_{n} \tag{1}
\end{align*}
$$

where $p, q, r$ and $a$ are functions on $V_{n-2}$. From Lemma 1 , it is easy to verify

Lemma 2 Let $\xi$ be the sequence of a bent function, $f$, on $V_{n}(n>2)$. Then $\xi$ is bent-based (linear-based) if and only if a $\left(x_{1}, \ldots, x_{n-2}\right)$ is the constant 1 (constant 0 ) in the expression for $f$ in (1).

Theorem 1 The conjecture of Adams and Tavares is true if only if every bent function is quadratic.

Proof. Let $\xi$ be the sequence of an arbitrary bent function on $V_{n}(n>2)$, say $f$.
Suppose any bent function is quadratic. It is easy to see that $a\left(x_{1}, \ldots, x_{n-2}\right)$ is constant in the expression for $f$ as in (1). By Lemma 2, the conjecture is true.

Conversely, suppose the conjecture is true i.e. $\xi$ is either bent-based or linear based. By Lemma $2, a\left(x_{1}, \ldots, x_{n-2}\right)$ is constant in the expression for $f$ as in (1). Suppose $f$ is not
quadratic. Then there exist two distinct indices $i$ and $j$ such that the coefficient of $x_{i} x_{j}$ in the expression for $f$ is not constant. Rewrite $f\left(x_{1}, \ldots, x_{n}\right)=g\left(x_{j_{1}}, \ldots, x_{j_{n-2}}, x_{i}, x_{j}\right)$, where $j_{1}, \ldots, j_{n-2}$ is any permutation of $\{1, \ldots, n\}-\{i, j\}$. Applying the conjecture to $g$ leads a contradiction.

## References

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